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# Design of LDPC Codes for Slepian-Wolf coding with uncertain knowledge of the correlation

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- Theoretical performance [SW73]
- Coding schemes based on channel codes [LXG02, LXG03, MUM10, XLC04]

In general, perfect knowledge of P(X, Y)

- P(X) unknown [JVW10]
- P(Y|X) given at decoder, partly unknown at encoder [Sga77]

Practical solution

• Feedback channel [AZG02, EY05, VAG06]

Context	Source Definition	Coding scheme	Experimental results	Conclusions
Contout				

#### In this work

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- P(Y|X) partly unknown to both encoder and decoder
- No feedback

### Objectives

- Design efficient coding/decoding schemes robust to uncertainty on P(Y|X)
- Solution based on non-binary LDPC codes











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### 2 Source Definition

### 3 Coding scheme

- 4 Experimental results
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# Modeling the uncertainty

Four source models considered in [DRK12]. Here, focus on the Static without Prior Source (SwP-Source)

#### Definition (SwP-Source)

A SwP-Source (X, Y), produces a sequence of independent discrete symbols  $\{(X_n, Y_n)\}_{n=1}^{+\infty}$  drawn from a distribution belonging to

 $\left\{ P(X, Y|\theta) = P(X)P(Y|X, \theta) \right\}_{\theta \in \mathscr{P}_{\theta}}$ 

 $\theta$  is fixed for  $\{(X_n, Y_n)\}_{n=1}^{+\infty}$ .

Context



3 Coding scheme

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### 5 Conclusions

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# LDPC-based Encoder in GF(q)

Theoretical Performance [Csi82]

$$R = \sup_{\theta \in \mathscr{P}_{\theta}} H(X|Y,\theta).$$





Degree distributions  $(\lambda(x), \rho(x))$  for H dimensioned for the worst case, need to be optimized

Two-step EM-based approach:  $(\mathbf{x}^{(\ell)}, \boldsymbol{\theta}^{(\ell)})$  at iteration  $\ell$ 

• LDPC decoding with estimate  $heta^{(\ell)}$  to obtain

 $P(X_n = k | y_n, \mathbf{s}, \boldsymbol{\theta}^{(\ell)})$ 

• Update of the  $\theta^{(\ell)}$  by maximizing

$$Q(\theta, \theta^{(\ell)}) = E_{\mathbf{X}|\mathbf{y}, \mathbf{s}, \theta^{(\ell)}} [\log P(\mathbf{y}|\mathbf{X}, \mathbf{s}, \theta)]$$
$$= \sum_{n=1}^{N} \sum_{k=0}^{q-1} P(X_n = k|y_n, \mathbf{s}, \theta^{(\ell)}) \log P(y_n|X_n = k, \theta)$$

#### Two-step EM-based approach

• LDPC decoding with estimate  $heta^{(\ell)}$  to obtain

 $P(X_n = k | y_n, \mathbf{s}, \boldsymbol{\theta}^{(\ell)})$ 

• Update of the  $\theta^{(\ell)}$  for  $Y = X \oplus Z$ ,  $P(Z = k) = \theta_k$ 

$$\theta_k^{(\ell+1)} = \frac{\sum\limits_{n=1}^N P(X_n = y_n \ominus k | y_n, \mathbf{s}, \theta^{(\ell)})}{\sum\limits_{n=1}^N \sum\limits_{k'=0}^{q-1} P(X_n = y_n \ominus k' | y_n, \mathbf{s}, \theta^{(\ell)})}$$

# Initialization of the EM algorithm (Additive Model)

Additive model:  $Y = X \oplus Z$ ,  $P(Z = k) = \theta_k$ 

• Compute

$$\mathbf{u} = \mathbf{s} \ominus H^T \mathbf{y} = H^T \mathbf{z}$$

Assumption: the  $U_m$  are obtained from *i.i.d.* R.V.s  $Z_i^{(m)}$ .

• Maximize

$$L(\theta) = \log P(\mathbf{u}|\theta) = \sum_{m=1}^{M} \log \mathscr{F}_{u_m}^{-1} \left( \prod_{j=1}^{dc} \mathscr{F}(W[h_j^{(m)}]\theta) \right)$$





### 3 Coding scheme





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Experimental results

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# Experimental framework

Symbols in GF(4), 
$$Y = X \oplus Z$$
,  
 $\theta = [\theta_0, ..., \theta_3]$  and  $Pr(Z = k) = \theta_k$ .  
 $\mathscr{P}_{\theta}$  s.t.  $\forall \theta \in \mathscr{P}_{\theta}, \ \theta_0 \ge 0.76$ .

Code tuned for the worst case  $\theta = [0.76, 0.08, 0.08, 0.08]$ 

• 
$$\lambda(x) = 0.413x + 0.375x^2 + 0.012x^4$$

- $\rho(x) = x$
- R = 1.6 bit/symbol

Context

Experimental results

Conclusions

# Initialization of the EM algorithm

MSE of the estimators



## Global scheme

1000 source vectors of length 1000 are generated. For each vector,  $\theta$  selected uniformly at random in  $\mathscr{P}_{\theta}$ .

Setup	Err	Time (s)	Rate (bit/symb.)
Genie-aided	$< 10^{-5}$	5.4	1.6
Learn. Seq.	$< 10^{-5}$	4.2	1.7
EM	$< 10^{-5}$	9.1	1.6
EM random	$7.2 imes10^{-3}$	47.0	1.6

# Conclusions

#### Summary

• Practical coding scheme based on non-binary LDPC codes when the correlation is uncertain

#### Future works

- Extension to the lossy case
- Design of good degree distributions
- Correlation model selection

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